

# Categorical Logic

Spring 2021

Homework 1 – Due March 3

## Algebraic Theories

1. (The theory with one sort) Let  $T_1$  be the signature with one sort, no function symbols and no axioms.
  - (a) Describe the set  $T(x_1, \dots, x_n)(\Sigma)$  of all terms in the context  $\vec{x} = (x_1, \dots, x_n)$ .
  - (b) Describe the morphisms  $(x_1, \dots, x_n) \rightarrow (y_1, \dots, y_m)$  in the syntactic category and prove that they are in bijection with functions  $\{1, 2, \dots, m\} \rightarrow \{1, 2, \dots, n\}$ .
  - (c) Let  $\mathcal{F}in$  be the category of finite sets. It is easy to see that  $\mathcal{F}in$  has finite sums. In consequence, its opposite category  $\mathcal{F}in^{op}$  has finite products.  
Construct an equivalence of categories  $\mathcal{S}yn(\Sigma) \simeq \mathcal{F}in^{op}$ . (Hint: use the fully faithful + essentially surjective characterization of equivalences)
  - (d) Describe the isomorphisms in  $\mathcal{S}yn(\Sigma)$ .
  - (e) A  $\Sigma$ -structure  $M$  is just a set. Find a morphism in  $\mathcal{S}yn(\Sigma)$  whose interpretation is the diagonal  $\Delta : M \rightarrow M \times M$ .

Remark: More generally, if  $\Sigma$  is the signature with set of sorts  $S$  and no function symbols, one can construct an equivalence  $\mathcal{S}yn(\Sigma) \simeq (\mathcal{F}in^S)^{op}$  (where  $\mathcal{F}in^S$  is the category of functors  $S \rightarrow \mathcal{F}in$ , i.e. the category of  $S$ -families of finite sets). It is a good exercise to try to prove it.

2. (Theories without function symbols) Let  $\Sigma$  be a signature with set of sorts  $S$  and no function symbols.
  - (a) Describe all the algebraic theories with signature  $\Sigma$  (in other words, list all the possible axioms that can be given).
  - (b) How many algebraic theory with signature  $\Sigma$  are there when  $S$  is a singleton ?
3. (Universal property of the theory with one sort) Let  $T_1$  be the algebraic theory of  $??$ . Let  $T = (\Sigma, A)$  another algebraic theory, and  $X$  a distinguished sort in  $\Sigma$ .
  - (a) Using  $X$ , construct a function  $T(T_1) \rightarrow T(T)$  between the set of terms.
  - (b) Prove that this function extends to a cartesian functor  $\mathcal{F}in^{op} = \mathcal{S}yn(T_1) \rightarrow \mathcal{S}yn(T)$  (no need to prove that is actually a functor, just explain what are the maps between the sets of objects and arrows, and justify that it is a cartesian functor).
4. (Equivalence of theories)
  - (a) Two theories are called "Lawvere equivalent" if they have equivalent syntactic categories. Consider the theory with two sorts  $X$  and  $Y$  and two function symbols  $f : X \rightarrow Y$  and  $g : Y \rightarrow X$  with axioms  $f(g(y)) = y$  and  $g(f(x)) = x$ . Prove that the syntactic category is equivalent to the one of the theory  $T_1$ .

- (b) Two theories are called "Morita equivalent" if they have equivalent categories of models in  $\text{Set}$ . Consider the theory  $T'$  with two sorts  $X$  and  $Y$  and two function symbols  $r : X \rightarrow Y$  and  $s : Y \rightarrow X$  with axioms  $r(s(y)) = y$  and the theory  $T''$ , with one sort  $X$  and one function symbol  $e : X \rightarrow X$  with axioms  $e(e(x)) = e(x)$ . Construct an equivalence between the categories of models in  $\text{Set}$  of both theories (Hint: use that, for a set  $X$  with an idempotent function  $e : X \rightarrow X$ , the set  $Y$  of the first theory is the set of fixed points of  $e$ ).
- (c) Prove that Lawvere equivalence  $\Rightarrow$  Morita equivalence, but not the contrary (use Question ?? for the second part)