

Many categories comes with distinguished structures
akin to algebraic operations (like $+$ \times ...)
on \mathbb{N}

ex Set - sum of 2 set $A+B$, empty set \emptyset $\emptyset \in A = A$
- product of 2 sets $A \times B$, singletons $\{x\}$ $\{x\} \times B = B$

other categories have also these structures
(Monoids, Cat ...)

initial objects

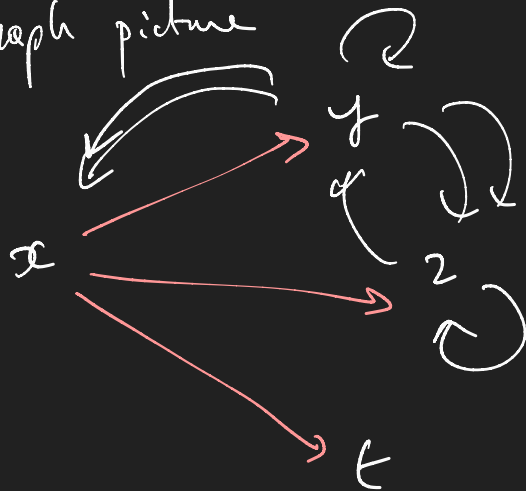
in \mathcal{C} cat. an object x is initial

if for any other object y in \mathcal{C}

$\text{Hom}_{\mathcal{C}}(x, y)$ is a singleton

(there is one morphism exactly
from x to y)

graph picture



Matrix picture

	y	z	t
y			
z	$\{*\}$	$\{*\}$	$\{*\} \dots$
x			

only singletons in row x

examples

- Set \emptyset the empty set is initial

$$\text{Hom}(\emptyset, E) = \{*\} \quad \forall E.$$

- Poset \emptyset the empty poset is initial

- Cat \emptyset the empty category is initial

- Mon (the cat of monoid)

$\{1\}$ monoid with one element

is initial.

$\{1\} \xrightarrow{\text{monoid morphism.}} M$

image of 1 has to be the unit of

$\mathbb{N} =$ free monoid on one generator.

$\{1\} =$ free monoid on zero generators.

- C is a poset.
an initial object corresponds
to a minimal element for the order.

• $\mathcal{P}(E) = \{\text{subset of } E\}$

minimal elt is $\emptyset \subset E$

is an initial object in $(\mathcal{P}(E), \subseteq)$
viewed as a category.

• (\mathbb{N}, \leq) 0 is initial

• (\mathbb{Z}, \leq) there is no / minimal elt
initial object

But not every category
admits an initial
object.

E set viewed as a cat.

has an initial object

iff $E = \{*\}$ is a
singleton.

Proposition initial objects are unique
up to unique isomorphism.

(any two initial objects are canonically isomorphic)

Proof: suppose x initial in \mathcal{C}

$\text{Hom}(x, x) = \{id_x\}$ is a singleton.

suppose x, x' two initial objects in \mathcal{C}

$\text{Hom}(x, x') = \{f\}$ $x \xrightarrow{f} x'$

$\text{Hom}(x', x) = \{g\}$ $x' \xrightarrow{g} x$

the composition $gf: x \rightarrow x$ is id_x

similarly $fg = id_{x'}$

so f and g are inverse to each other.

so x and x' are isomorphic.

and this isomorphism between them is unique.

□

ex all empty sets are isomorphic.



$\exists!$ bij. between empty sets

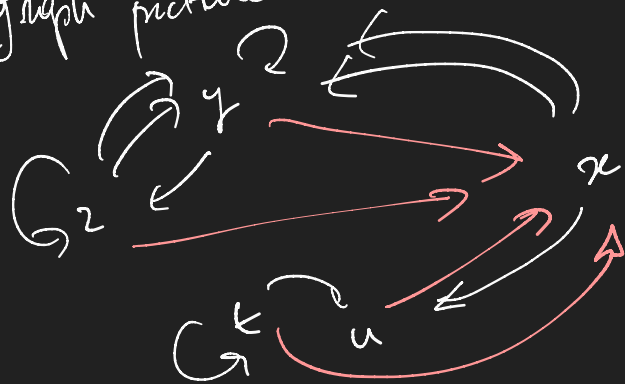
terminal object

in \mathcal{C} cat an object x is terminal
(or final)

if for any y in \mathcal{C}

$\text{Hom}_{\mathcal{C}}(y, x)$ is a singleton

graph picture



matrix picture

\rightarrow	y	x	\vdots
y			
x	$\text{Hom}(x, y)$		
\vdots			

Column of
singletons.

Recall C^{op} the opposite of C

preposition if x is initial
(or terminal) in C , then
 x^{op} is terminal (or initial)
in C .

The two definitions of initial/terminal
object are "dual" to each
(exchanged by passing to the
opposite cat)

examples

- Set any singleton
is terminal.

- Poset $\{1\}$ poset with
one elt
is terminal.

- Cat the punctual cat
is terminal.

- Monoids. $\{1\}$ ^{group} Monoid
^{Groups} with
single elt
is also terminal

- in a poset P (viewed as a cat.)
an object is terminal iff it is
a maximal element

• $P(E) = \{ \text{subsets of } E \}$
 E is a terminal object

• (\mathbb{N}, \leq) no maximal elt.
no terminal object.

a category C is
called pointed

if it has an initial
object 0 , a terminal
object 1 , such that

$$0 \xrightarrow{\sim} 1$$

is an isomorphism.

ex Groups, Monoids,

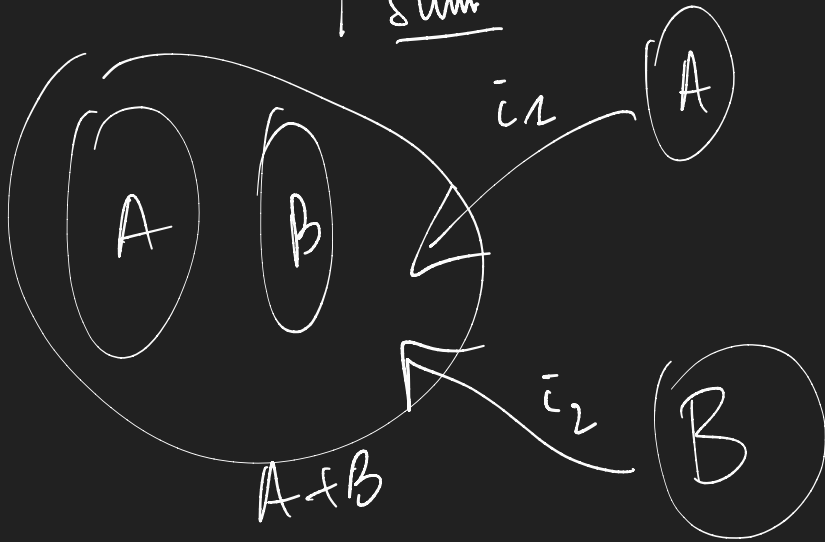
Category with sums

in $\text{Set} \Rightarrow A, B$

distinguished set $A+B$ ($A \perp B$)

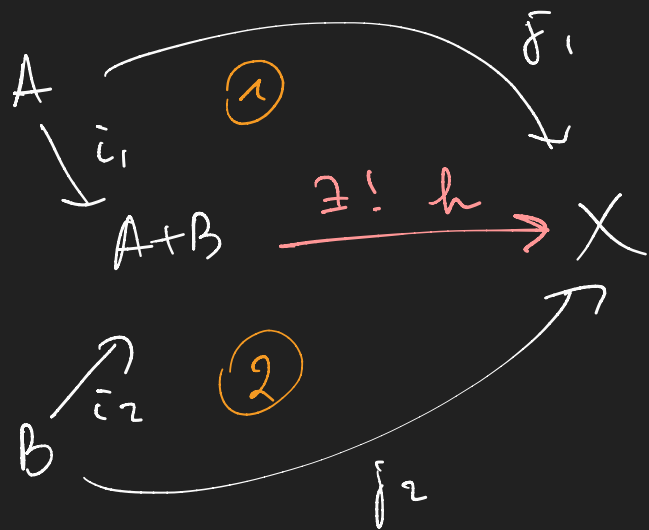
called the disjoint union of A and B

sum



$A+B$ comes equipped with distinguished inclusions

$$A \xrightarrow{i_1} A+B \xleftarrow{i_2} B$$



there exists a unique
map $A+B \xrightarrow{h} X$
such that

$$h \circ i_1 = j_1$$

$$h \circ i_2 = j_2$$

universal property of $A+B$

the triangle (1)
commutes

the triangle (2)
commutes.

the whole
diagram commutes.

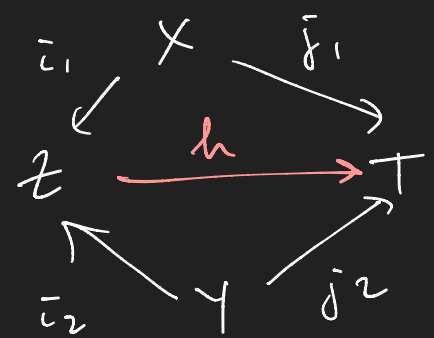
in a category \mathcal{C} , given two objects X, Y

the sum of X and Y is an object Z
coproduct

together with maps $X \xrightarrow{i_1} Z \xleftarrow{i_2} Y$

such that, for any other diagram $X \xrightarrow{j_1} T \xleftarrow{j_2} Y$,

there exists a unique map $Z \xrightarrow{h} T$ such that the diagram



commutes.

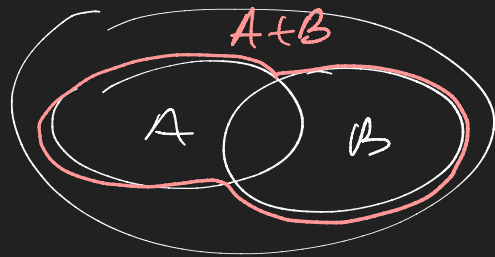
examples

- in Set, get back the usual notion of disjoint union
sum
coproducts
- in Poset, get the usual sum
- in Cat, get the sum
- in a poset P viewed as a category this is the notion of join of two elements.

• $P(E)$

$A \subseteq E$ $B \subseteq E$

$A + B$ is $A \cup B$ the union of A and B in E



• in (\mathbb{N}, \leq)

the sum of two element is

the max (and not the arithmetic sum)

- Monoids have sums.

this is the notion of "free product"
of two monoids.

$\text{List}(E)$ = free monoid on the set E

$$\underbrace{\text{List}(E) + \text{List}(E')}_{\text{sum in the cat. of monoids.}} = \text{List}(\underbrace{E + E'}_{\substack{\text{sum in} \\ \text{the cat. of} \\ \text{Sets.}}})$$

(exo: try to prove this using the univ. prop. of the free monoid)

- E a set
viewed as a cat.
does not have all
sums.

if $x \neq y$

$x + y$ does not
exist

but

$$x + x = x$$

(exercise)

\mathcal{C} cat. X, Y objects

proposition if the sum of X and Y exists
it is unique up to unique isomorphism.

proof let Z be a sum for X and Y .

$$X \xrightarrow{\bar{c}_1} Z \xleftarrow{\bar{c}_2} Y$$

there is a unique morphism $Z \xrightarrow{1_Z} Z$

s.t.

$$\begin{array}{ccc} & X & \\ \bar{c}_1 \swarrow & & \searrow \bar{c}_1 \\ Z & \xrightarrow{1_Z} & Z \\ \bar{c}_2 \swarrow & & \searrow \bar{c}_2 \\ & Y & \end{array}$$

commutes.

has to be the
identity of Z !

assume $x \xrightarrow{i_1} z \xleftarrow{i_2} y$
 and $x \xrightarrow{j_1} z' \xleftarrow{j_2} y$

are two sums.

by prop. of z , there exists
 a unique map $z \xrightarrow{f} z'$

s.t.
$$\begin{array}{ccc} i_1 & \searrow & \\ x & & \\ j_1 & \swarrow & \\ z & \xrightarrow{f} & z' \\ i_2 & \swarrow & \\ y & & \\ j_2 & \searrow & \end{array}$$
 commutes.

by prop of z' , there exists
 a unique map $z' \xrightarrow{g} z$ s.t.

$$\begin{array}{ccc} & x & \\ j_1 \swarrow & & \searrow i_1 \\ & z' & \xrightarrow{g} z \\ & & \swarrow i_2 \\ j_2 \swarrow & & \nearrow y \end{array}$$
 commutes.

hence $fg = id_z$, $gf = id_{z'}$

so f, g are inverse of each other
 and z and z' are isomorphic.

uniquely

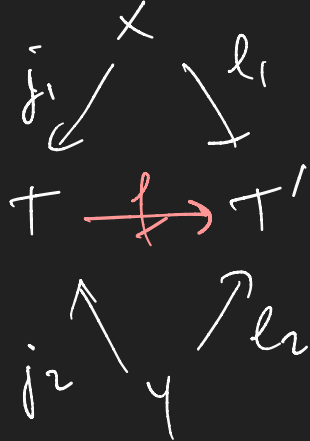
It is possible to view sums as initial objects

define a category cocones (X, Y) $(X, Y \text{ fixed in } \mathcal{C})$

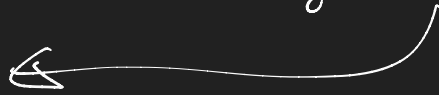
objects : diagrams $X \xrightarrow{j_1} T \xleftarrow{j_2} Y$

a morphism from $X \xrightarrow{j_1} T \xleftarrow{j_2} Y$ to $X \xrightarrow{l_1} T' \xleftarrow{l_2} Y$

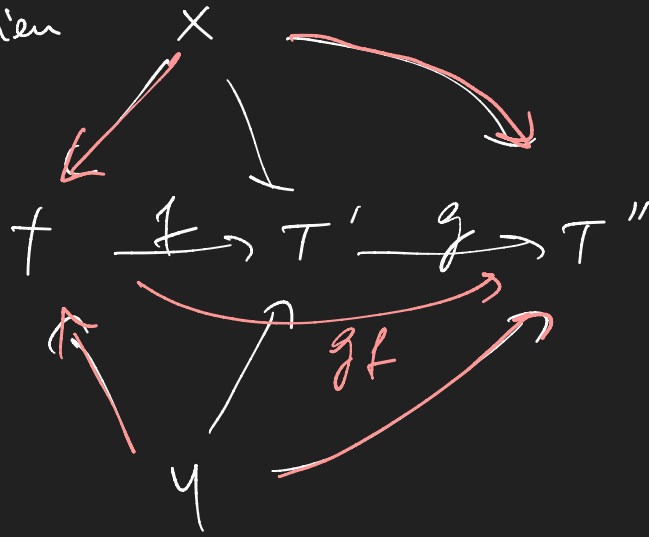
is a diagram



(equivalently a morphism is
a map $T \xrightarrow{f} T'$ s.t
the diagram commutes)



composition



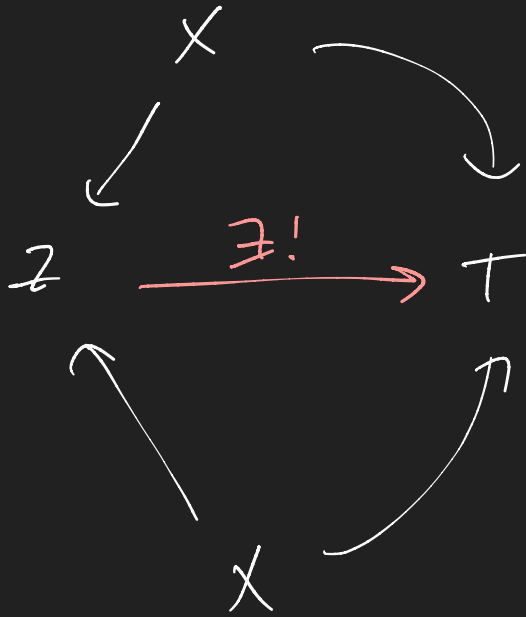
ex check this is a cat.

proposition

an initial object in
locome (X, Y)

is the same thing as
a sum for X and Y .

proof (just unravels
the definition of
an initial object)



initial in

$\text{closure}(X, Y)$