

Where are we?

CoT theory is a synthesis for many structures. ✓

classification tool ✓

language ← here

calculus (later)

the "Hom" functor (the most important functor of all)

\mathcal{C} a category $x \in \text{ob}(\mathcal{C})$ an object

$$\text{Hom}_{\mathcal{C}}(x, -) : \mathcal{C} \longrightarrow \text{Set}$$

$$\begin{array}{ccc} y & \xrightarrow{\quad} & \text{Hom}_{\mathcal{C}}(x, y) \\ f \downarrow & & \vdots \text{Hom}(x, f) \\ z & \xrightarrow{\quad} & \text{Hom}_{\mathcal{C}}(x, z) \end{array}$$

$$\begin{array}{ccc} x & \xrightarrow{g} & y \\ \downarrow & & \vdots f \\ x & \xrightarrow{fg} & z \end{array}$$

(ex: check this is a functor)

example : $\mathcal{C} = \text{Set}$ E a set

$$\text{Hom}(E, -) : \text{Set} \longrightarrow \text{Set}$$
$$A \longmapsto \text{Hom}(E, A) = A^E$$

if $E = \{0, 1\}$ has two elt. $\text{Hom}(E, A) = A^2$

the "other" Hom functor

$$\mathcal{C} \text{ cat} \quad x \in \text{ob}(\mathcal{C})$$

$$\text{Hom}(-, x) : \mathcal{C} \longrightarrow \text{Set}$$

$$y \longmapsto \text{Hom}(y, x)$$

$$f \downarrow$$

$$\times$$

$$\uparrow \text{Hom}(f, x)$$

$$f \downarrow$$

$$z \longmapsto \text{Hom}(z, x)$$

$$z \xrightarrow{h} x$$

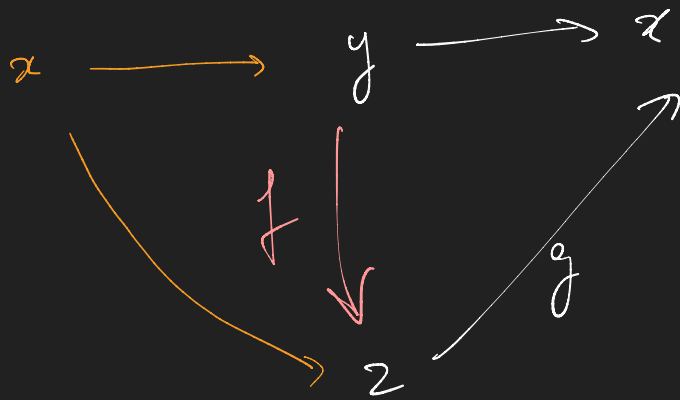
going the opposite way of f

$$y \xrightarrow{g} x$$

strictly speaking this is not a functor

even if send identities to identities and respect composition.

$$\text{Hom}(x, y) \xrightarrow{f \circ -} \text{Hom}(x, z)$$



$$\begin{array}{ccc} gf & \text{Hom}(y, x) & \\ \uparrow & \uparrow - \circ f & \\ g & \text{Hom}(z, x) & \end{array}$$

This motivates the definition of contravariant functors

$f: C \xrightarrow{\text{contra}} D$ is the data

$$f_0: \text{ob}(C) \rightarrow \text{ob}(D)$$

$$f_1: \text{Arr}(C) \rightarrow \text{Arr}(D)$$

such that for any $u: x \rightarrow y$ in C

$$\text{dom}(f_1(u)) = f_0(\text{cod}(u))$$

$$\text{cod}(f_1(u)) = f_0(\text{dom}(u))$$

f sends id to id .
 f respects compositions

$$\begin{array}{ccc} C & \xrightarrow{\text{contra}} & D \\ x & \longmapsto & f_0(x) \\ u \downarrow & & \uparrow f_1(u) \\ y & \longmapsto & f_0(y) \end{array}$$

in CT. there is two kinds of functors
the usual ones (called "covariant")

and the contravariant ones

not so convenient to have two kinds of functors

There's a trick to describe a contravariant functor

by means of a covariant one :

use the opposite categories

the opposite of a category

C is a cat its opposite is denoted C^{op} is the cat.

- $ob(C^{op}) = ob(C)$ (same objects)

the opposite of an object x is denoted x^{op}

- $Hom_{C^{op}}(x^{op}, y^{op}) = Hom_C(y, x)$

(in other terms $x^{op} \xrightarrow{f^{op}} y^{op}$ is the same thing as $y \xrightarrow{f} x$)

- identity arrows: $1_{x^{op}} = (1_x)^{op}$

- composition: $x^{op} \xrightarrow{f^{op}} y^{op} \xrightarrow{g^{op}} z^{op} \iff x \xleftarrow{f} y \xleftarrow{g} z$

$g^{op} \cdot f^{op} = (fg)^{op}$

proposition

a contravariant functor $\mathcal{C} \xrightarrow{\text{contra}} \mathcal{D}$ is

the same thing as a covariant functor $\mathcal{C}^{\text{op}} \rightarrow \mathcal{D}$

ex : ~~$\text{Hom}_{\mathcal{C}}(-, x) : \mathcal{C} \xrightarrow{\text{contra}} \text{Set}$~~

$\text{Hom}_{\mathcal{C}}(-, x) : \mathcal{C}^{\text{op}} \xrightarrow{=} \text{Set}$

(the most important
contravariant functor)

example

$\mathcal{C} = \text{Set}$

$\text{Set}^{\text{op}} \longrightarrow \text{Set}$

$E \longmapsto \mathcal{P}(E) = \text{set of subsets of } E$
 $= \text{Hom}_{\text{Set}}(E, \{0,1\})$

= example of a
contravariant Hom functor.

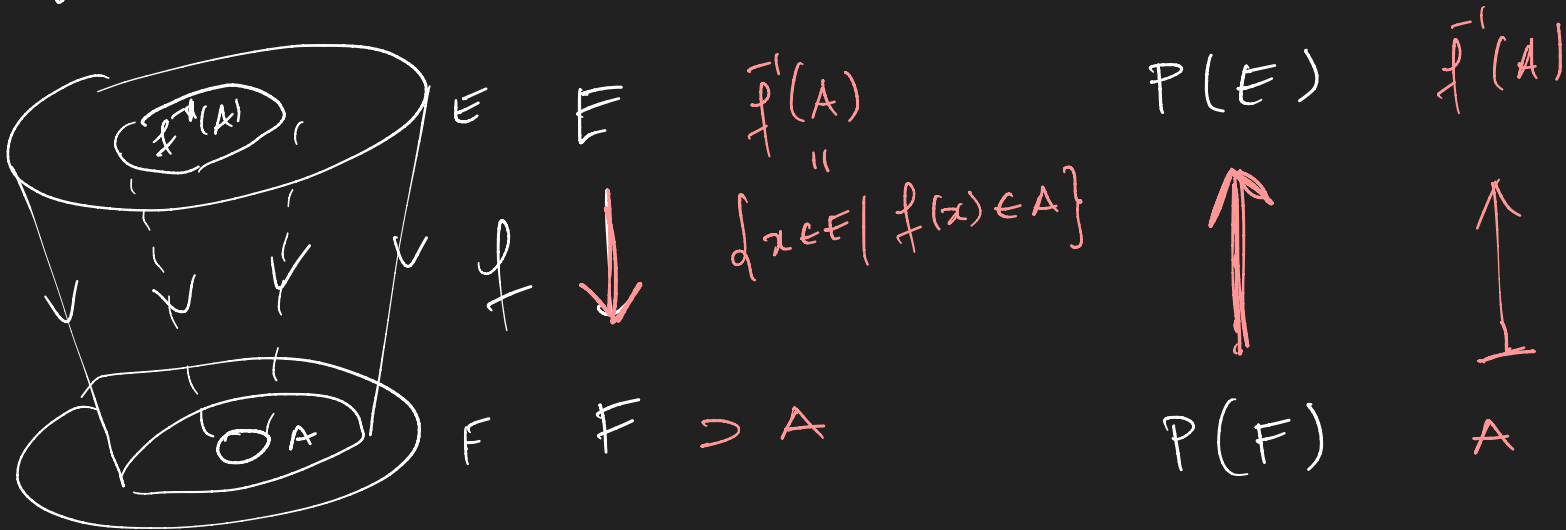
bijection $\mathcal{P}(E) \cong \text{Hom}(E, \{0,1\})$

$A \subseteq E \mapsto$ characteristic
function

$E \rightarrow \{0,1\}$

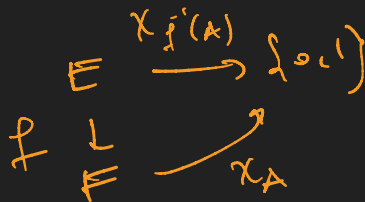
$x \mapsto \begin{cases} 1 & \text{if } x \in A \\ 0 & \text{if not} \end{cases}$

Subsets and functions

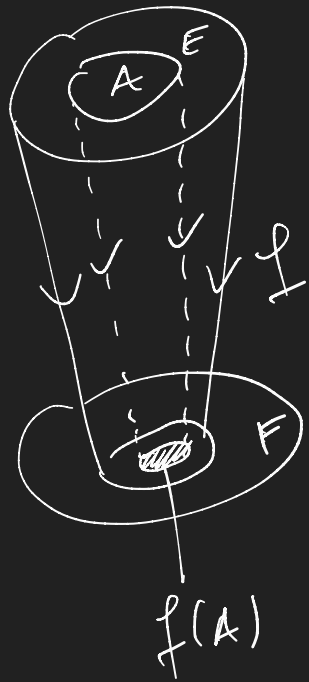


exercise: check it.

char. fct of $f^{-1}(A)$ = composition of char fct of A with f .



Remark: there is another way to transport subsets along functions $f: E \rightarrow F$: the image of a subset



$$\begin{array}{ccc}
 E \supset A & & \mathcal{P}(E) \\
 \downarrow & & \downarrow \\
 F \supset \tilde{f}(A) & & \mathcal{P}(F) \\
 & & \text{"}
 \end{array}$$

$$\left\{ y \in F \mid \exists \underline{x} \in A \text{ st. } f(x) = y \right\}$$

image of A by f

the association of $P(E)$ to E defines two functors,

a contravariant one

$$\begin{array}{ccc} \text{Set}^{\text{op}} & \longrightarrow & \text{Set} \\ E & \longmapsto & P(E) \\ \downarrow f & & \uparrow f^{-1} \\ F & \longmapsto & P(F) \end{array} \quad \begin{array}{l} \text{invert} \\ \text{image} \end{array}$$

a covariant one

$$\begin{array}{ccc} \text{Set} & \longrightarrow & \text{Set} \\ E & \longmapsto & P(E) \\ \downarrow f & & \downarrow \tilde{f} \\ F & \longmapsto & P(F) \end{array} \quad \begin{array}{l} \text{direct} \\ \text{image} \end{array}$$

a functor is not determined by what it does on objects only
(these two functors have the same action on objects)
but are very different.

the opposition is a way to produce a cat. from a cat.

other ways to produce a cat from some categories =

- the sum of categories. C, D are 2 cat.

$C + D$ ($C \perp\!\!\!\perp D$) is the cat. s.t.

$$\text{ob}(C + D) = \text{ob}(C) + \text{ob}(D)$$

$$\text{Arr}(C + D) = \text{Arr}(C) + \text{Arr}(D)$$

(no arrow from an object of C to an object of D
and vice-versa)

id, comp. are defined by those of C and D

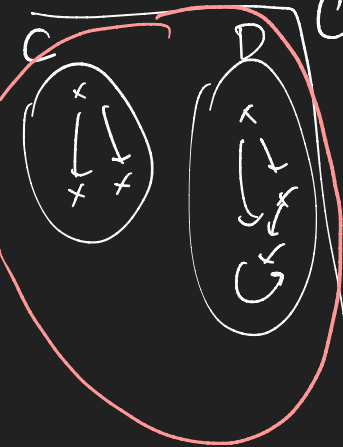
identity
arrows

if $x \in C$

use $1_x^{\text{in } C}$

if $x \in D$

use $1_x^{\text{in } D}$



$C + D$

- the product of categories C, D 2 cat

$C \times D$ (the product of C and D) is the cat:

$$\begin{aligned} \bullet \text{ ob}(C \times D) &= \text{ob}(C) \times \text{ob}(D) \\ &= \text{pairs } (x, y) \quad \begin{array}{l} x \in \text{ob}(C) \\ y \in \text{ob}(D) \end{array} \end{aligned}$$

$$\bullet \text{ Arr}(C \times D) = \text{Arr}(C) \times \text{Arr}(D)$$

$$\begin{array}{ccc} (x, y) & \longrightarrow & (x', y') \\ \left(\begin{array}{c} x \xrightarrow{f} x' \\ \text{in} \\ C \end{array} , \begin{array}{c} y \xrightarrow{g} y' \\ \text{in} \\ D \end{array} \right) & & \end{array}$$

• identity of (x, y)

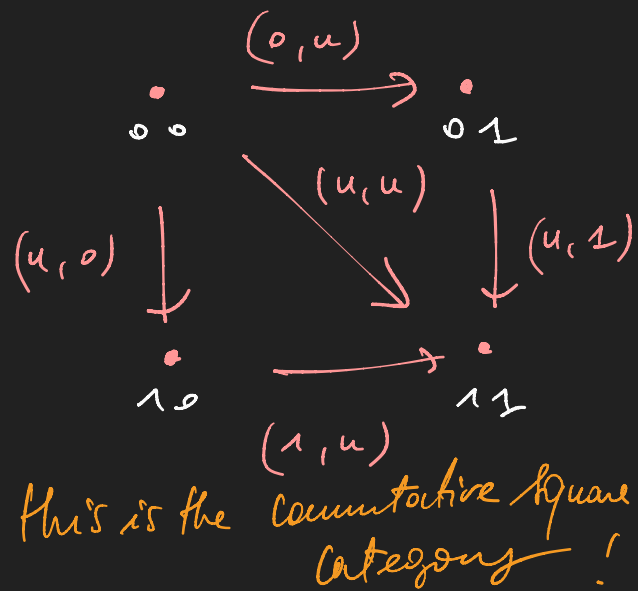
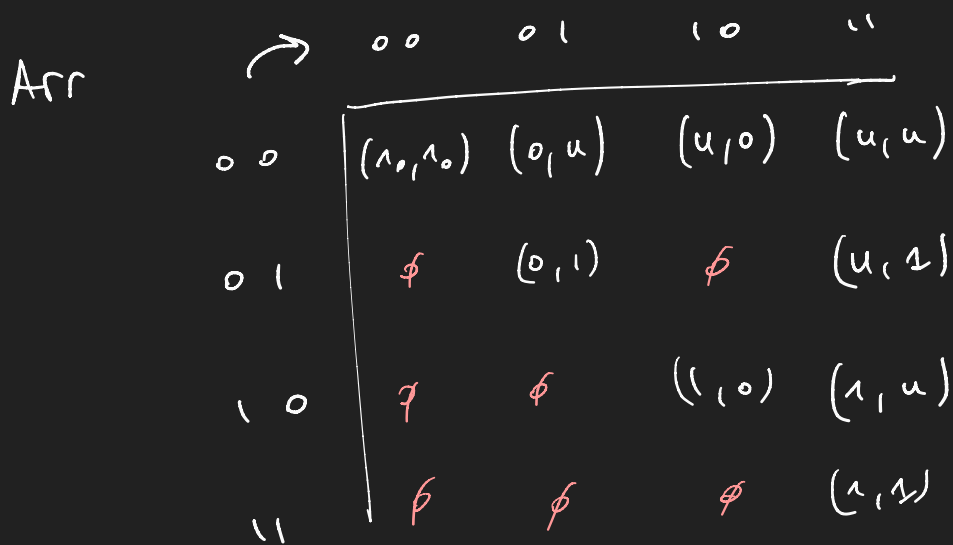
is $(1_x, 1_y)$

• composition
as expected..

example $[1] = \{ 0 \xrightarrow{u} 1 \}$

$ob([1]) = \{ 0, 1 \}$ $Arr([1]) = \{ \overset{0}{1}_0, \overset{1}{1}_1, u \}$

$[1] \times [1]$ $ob = \{ 0, 1 \} \times \{ 0, 1 \} = \{ (0,0), (0,1), (1,0), (1,1) \}$



example $\mathbb{1} = \{*\}$ the punctual category

\mathcal{C} is any cat.

$$\mathbb{1} \times \mathcal{C} = \mathcal{C}$$

(isomorphic to \mathcal{C} .)

for the sum $\emptyset =$ the empty category

$$\emptyset + \mathcal{C} = \mathcal{C}$$

(isomorphic categories)

we saw two Hom functors C cat, $x \in C$
 $y \in C$

$$\text{Hom}_C(x, -) : C \longrightarrow \text{Set}$$

$$\text{Hom}_C(-, y) : C^{\text{op}} \longrightarrow \text{Set}$$

can be put together in a single functor:

$$C^{\text{op}} \times C \longrightarrow \text{Set}$$

$$(x^{\text{op}}, y) \longmapsto \text{Hom}_C(x, y)$$

If C_1, C_2, D are categories.

a functor of two variables is a functor

$$C_1 \times C_2 \longrightarrow D$$

ex $\text{Hom}(-, -) : C^{\text{op}} \times C \longrightarrow \text{Set}$

is a functor of 2 variables. which is
contravariant in the first variable
and covariant in the second var.